

# Regular and Context-Free Pattern Languages Over Small Alphabets

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DLT 2012

# Pattern languages 1/2

$\Sigma$ : Finite alphabet of terminal symbols (e. g.  $\Sigma = \{a, b, c, d\}$ ).

$X = \{x_1, x_2, x_3, \dots\}$ : Infinite alphabet of variables.

A word  $\alpha \in (\Sigma \cup X)^+$  is called a *pattern*.

# Basic Definitions

*Morphism*: Mapping  $h : \Gamma_1^* \rightarrow \Gamma_2^*$  with  $h(x \cdot y) = h(x) \cdot h(y)$ ;  
 $h$  is *nonerasing* iff, for every  $a \in \Gamma_1$ ,  $h(a) \neq \varepsilon$ .

*Substitution*: Morphism  $h : (\Sigma \cup X)^* \rightarrow \Sigma^*$  with  $h(a) = a$  for all  
 $a \in \Sigma$ .

$L_{E,\Sigma}(\alpha) := \{h(\alpha) \mid h \text{ is a substitution}\}$ .

$L_{NE,\Sigma}(\alpha) := \{h(\alpha) \mid h \text{ is nonerasing substitution}\}$ .

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$$L_{NE, \{a,b,c\}}(\alpha) = \{w \mid w = u \text{ aa } v \text{ } u \text{ } v \text{ cb } u, u, v \in \{a, b, c\}^+\}.$$

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$$ccbaaccbcbccb \notin L_{NE, \{a,b,c\}}(\alpha)$$

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# Notation

- $\text{var}(\alpha)$ : Set of variables occurring in  $\alpha$ .  
E. g.  $\text{var}(x_1 abx_2 bax_1 x_2 cx_3) = \{x_1, x_2, x_3\}$ .
- $|\alpha|_{x_i}$ : Number of occurrences of variable  $x_i$  in  $\alpha$ .  
E. g.  $|x_1 abx_2 bax_1 x_2 cx_3|_{x_2} = 2$ .
- REG := regular languages,
- CF := context-free languages,
- CS := context-sensitive languages,
- $Z \in \{E, NE\}$ , Z-PAT $_{\Sigma}$  = Z-Pattern Languages (w. r. t.  $\Sigma$ ).

# Pattern Languages and the Chomsky Hierachy

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  - $Z\text{-PAT}_\Sigma \not\subseteq REG$ ,  $Z\text{-PAT}_\Sigma \not\subseteq CF$ ,
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  - $Z\text{-PAT}_\Sigma \subseteq CS$ ,
  - $Z\text{-PAT}_\Sigma \# REG$ ,  $Z\text{-PAT}_\Sigma \# CF$ ,
  - $Z\text{-PAT}_\Sigma \cap REG \neq \emptyset$ ,  $Z\text{-PAT}_\Sigma \cap CF \neq \emptyset$ .
- The sets  $Z\text{-PAT}_\Sigma \cap REG$  and  $Z\text{-PAT}_\Sigma \cap CF$  are surprisingly difficult to characterise.

# Notation

$$L_{E,\Sigma}(\alpha) \in \text{REG} \Leftrightarrow \alpha \in \text{REG} (E, \Sigma),$$
$$L_{Z,\Sigma}(\alpha) \in \text{REG}, \text{ for every } Z \in \{E, \text{NE}\} \Leftrightarrow \alpha \in \text{REG} (Z, \Sigma).$$

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$x_1 a x_2 b x_3 c x_4 \in \text{REG} (\text{NE}, \{a, b, c, d\}),$

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- $a x_1 a x_2 a x_1 a x_3 \in \text{CF} \setminus \text{REG}, (\text{Z}, \{a, b\}),$
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$x_1 a x_2 x_3^2 b x_4 x_3 x_5 x_6 x_7^2 x_8 x_9 x_5 x_3 a x_4 x_5 x_4 x_9 x_{10} b x_{11} a x_{12} b x_{13} a x_{14} x_{15} b x_{15}^2 x_{16}^2 b x_{17}$   
 $\in \text{REG} (\text{E}, \{a, b\}).$



# Regular and Block-Regular Patterns 1/2

- *Regular patterns*

- A pattern is a *regular* pattern iff every variable has only one occurrence,
- e. g.,  $x_1 a x_2 b x_3 c x_4$ .
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- *Block-regular patterns*

- A pattern is a *block-regular* pattern iff every *variable block* contains a variable with only one occurrence in the whole pattern,
- e. g.,  $x_1 x_2 a x_1 x_3 x_4 x_1 x_4 b x_5 x_6 x_4 c x_7$ ,
- $Z \in \{E, NE\}$ ,  $Z\text{-PAT}_{\Sigma, \text{b-reg}} = Z\text{-Pattern Languages (w. r. t. } \Sigma)$  defined by block-regular patterns.

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- Hard cases:
  - NE case in general,
  - E case for alphabets of size 2 and 3.

# Repeated Variables

## Theorem (Jain et al., 2010)

Let  $|\Sigma| \geq 2$  and let  $\alpha \in X^+$ . If, for every  $x \in \text{var}(\alpha)$ ,  $|\alpha|_x \geq 2$ , then  $\alpha \notin \text{CF}(\mathbf{E}, \Sigma)$  and  $\alpha \notin \text{REG}(\mathbf{NE}, \Sigma)$ .



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## Theorem

Let  $|\Sigma| \geq 2$  and let  $\alpha \in (\Sigma \cup X)^+$ . If, for every  $x \in \text{var}(\alpha)$ ,  $|\alpha|_x \geq 2$ , then  $\alpha \notin \text{REG}(\mathbf{Z}, \Sigma)$ .

## Theorem

Let  $|\Sigma| \geq 3$  and let  $\alpha \in (\Sigma \cup X)^+$ . If, for every  $x \in \text{var}(\alpha)$ ,  $|\alpha|_x \geq 2$ , then  $\alpha \notin \text{CF}(\mathbf{Z}, \Sigma)$ .

# Spread Variable Blocks

## Theorem

Let  $\Sigma$  be terminal alphabets with  $|\Sigma| \leq 3$  and let  $Z \in \{E, NE\}$ . Let  $\alpha$  be a pattern with

$$\alpha = \beta \cdot d \cdot \gamma \cdot d' \cdot \delta, \text{ where } d, d' \in \Sigma, \gamma \in X^+.$$

If  $\text{var}(\gamma) \subseteq \text{var}(\beta) \cup \text{var}(\delta)$ , then  $\alpha \notin \text{REG}(Z, \Sigma')$ .

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# Example

E case:

$$|\Sigma| = 2 : x_1 a x_2 x_2 a x_3 \in \text{REG},$$

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E case:

$|\Sigma| = 2 : x_1 \mathbf{a} x_2 x_2 \mathbf{a} x_3 \in \text{REG},$

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E case:

$|\Sigma| = 2$  :  $x_1 \mathbf{a} x_2 x_2 \mathbf{a} x_3 \in \text{REG}, x_1 \mathbf{a}(\mathbf{bb})^* \mathbf{a} x_3$

$|\Sigma| \geq 3$  :  $x_1 \mathbf{a} x_2 x_2 \mathbf{a} x_3 \notin \text{REG},$

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$L_{E,\{a,b\}}(x_1 d_1 x_2 x_2 d_2 x_3 x_3 d_3 x_4) \notin \text{REG}$  iff  $d_1 = d_2 = d_3.$

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## Two Basic Lemmas

Lemma 1:  $L_{E,\{a,b\}}(\beta \cdot y \cdot \beta' \cdot a \cdot \gamma \cdot b \cdot \delta' \cdot z \cdot \delta) = L_{E,\{a,b\}}(\beta \cdot y \cdot ab \cdot z \cdot \delta)$ ,

- $\beta, \delta \in (\{a, b\} \cup X)^*$ ,
- $\beta', \gamma, \delta' \in X^*$ ,
- $y, z \in X, |\alpha|_y = |\alpha|_z = 1$ ,
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Lemma 2:  $L_{E,\{a,b\}}(\beta \cdot y \cdot \beta' \cdot a \cdot \gamma \cdot a \cdot \delta' \cdot z \cdot \delta) = L_{E,\{a,b\}}(\beta \cdot y \cdot a(bb)^* a \cdot z \cdot \delta)$ ,

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- $\text{var}(\beta' \cdot \gamma \cdot \delta') \cap \text{var}(\beta \cdot \delta) = \emptyset$ ,
- there exists a  $z' \in \text{var}(\gamma)$  with  $|\alpha|_{z'} = 2$ ,
- $|\gamma|_x$  is even,  $x \in \text{var}(\gamma)$ .

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$$x_1 a x_2 x_3^2 b x_4 x_3 x_5 x_6 x_7^2 x_8 x_9 x_5 x_3 a x_4 x_5 x_4 x_9 x_{10} b x_{11} a x_{12} b x_{13} a x_{14} x_{15} b x_{15}^2 x_{16}^2 b x_{17}$$

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$$\Rightarrow L_{E, \{a, b\}}(x_1 a x_2 b a x_{10} b x_{11} a x_{12} b x_{13} a x_{14} b (aa)^* b x_{17}) \in \text{REG.}$$

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...but we have to be careful with “copy-like” languages.